

## FORCING EDGE DETOUR NUMBER OF AN EDGE DETOUR GRAPH

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ABSTRACT. For two vertices  $u$  and  $v$  in a graph  $G = (V, E)$ , the *detour distance*  $D(u, v)$  is the length of a longest  $u$ - $v$  path in  $G$ . A  $u$ - $v$  path of length  $D(u, v)$  is called a  $u$ - $v$  *detour*. A set  $S \subseteq V$  is called an *edge detour set* if every edge in  $G$  lies on a detour joining a pair of vertices of  $S$ . The *edge detour number*  $dn_1(G)$  of  $G$  is the minimum order of its edge detour sets and any edge detour set of order  $dn_1(G)$  is an *edge detour basis* of  $G$ . A connected graph  $G$  is called an *edge detour graph* if it has an edge detour set. A subset  $T$  of an edge detour basis  $S$  of an edge detour graph  $G$  is called a *forcing subset* for  $S$  if  $S$  is the unique edge detour basis containing  $T$ . A forcing subset for  $S$  of minimum cardinality is a *minimum forcing subset* of  $S$ . The *forcing edge detour number*  $fdn_1(S)$  of  $S$ , is the minimum cardinality of a forcing subset for  $S$ . The *forcing edge detour number*  $fdn_1(G)$  of  $G$ , is  $\min\{fdn_1(S)\}$ , where the minimum is taken over all edge detour bases  $S$  in  $G$ . The general properties satisfied by these forcing subsets are discussed and the forcing edge detour numbers of certain classes of standard edge detour graphs are determined. The parameters  $dn_1(G)$  and  $fdn_1(G)$  satisfy the relation  $0 \leq fdn_1(G) \leq dn_1(G)$  and it is proved that for each pair  $a, b$  of integers with  $0 \leq a \leq b$  and  $b \geq 2$ , there is an edge detour graph  $G$  with  $fdn_1(G) = a$  and  $dn_1(G) = b$ .

*Key words:* detour, edge detour set, edge detour basis, edge detour number, forcing edge detour number.

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### 1. INTRODUCTION

By a *graph*  $G = (V, E)$ , we mean a finite undirected graph without loops or multiple edges. The order and size of  $G$  are denoted by  $p$  and  $q$  respectively. For basic definitions and terminologies, we refer to [1, 8].

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For vertices  $u$  and  $v$  in a connected graph  $G$ , the *detour distance*  $D(u, v)$  is the length of a longest  $u$ - $v$  path in  $G$ . A  $u$ - $v$  path of length  $D(u, v)$  is called a  $u$ - $v$  *detour*. It is known that the detour distance is a metric on the vertex set  $V$ . Detour distance and detour center of a graph were studied by Chartrand et al. in [2, 7].

A vertex  $x$  is said to lie on a  $u$ - $v$  detour  $P$  if  $x$  is a vertex of  $P$  including the vertices  $u$  and  $v$ . A set  $S \subseteq V$  is called a *detour set* if every vertex  $v$  in  $G$  lies on a detour joining a pair of vertices of  $S$ . The *detour number*  $dn(G)$  of  $G$  is the minimum order of a detour set and any detour set of order  $dn(G)$  is called a *detour basis* of  $G$ . A vertex  $v$  that belongs to every detour basis of  $G$  is a *detour vertex* in  $G$ . If  $G$  has a unique detour basis  $S$ , then every vertex in  $S$  is a detour vertex in  $G$ . These concepts were studied by Chartrand et al. in [3] and have interesting applications in Channel Assignment Problem in radio technologies [4, 9].

An edge  $e$  of  $G$  is said to lie on a  $u$ - $v$  detour  $P$  if  $e$  is an edge of the path  $P$ . In general, there are graphs  $G$  for which there exist edges which do not lie on a detour joining any pair of vertices of  $V$ . For the graph  $G$  given in Figure 1, the edge  $v_1v_2$  does not lie on a detour joining any pair of vertices of  $V$ . This motivated us to introduce the concepts of *weak edge detour set* of a graph [10] and *edge detour graphs* [11].

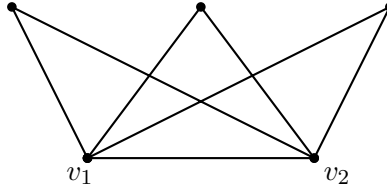


FIGURE 1.  $G$

A set  $S \subseteq V$  is called a *weak edge detour set* of  $G$  if every edge in  $G$  has both its ends in  $S$  or it lies on a detour joining a pair of vertices of  $S$ . The *weak edge detour number*  $dn_w(G)$  of  $G$  is the minimum order of its weak edge detour sets and any weak edge detour set of order  $dn_w(G)$  is called a *weak edge detour basis* of  $G$ . Weak edge detour sets and weak edge detour number of a graph were introduced and studied by Santhakumaran and Athisayanathan in [10].

A set  $S \subseteq V$  is called an *edge detour set* of  $G$  if every edge in  $G$  lies on a detour joining a pair of vertices of  $S$ . The *edge detour number*  $dn_1(G)$  of  $G$  is the minimum order of its edge detour sets and any edge detour set of order  $dn_1(G)$  is an *edge detour basis* of  $G$ . A graph  $G$  is called an *edge detour graph* if it has an edge detour set. A vertex  $v$  in an edge detour graph  $G$  is an *edge detour vertex* if  $v$  belongs to every edge detour basis of  $G$ . If  $G$  has a unique edge detour basis  $S$ , then every vertex in  $S$  is an edge detour vertex

of  $G$ . Edge detour graphs were introduced in [11] and further studied in [12] by Santhakumaran and Athisayanathan.

For the graph  $G$  given in Figure 2 (a), the sets  $S_1 = \{u, x\}$ ,  $S_2 = \{u, w, x\}$  and  $S_3 = \{u, v, x, y\}$  are detour basis, weak edge detour basis and edge detour basis of  $G$  respectively and hence  $dn(G) = 2$ ,  $dn_w(G) = 3$  and  $dn_1(G) = 4$ . For the graph  $G$  given in Figure 2 (b), the set  $S = \{u_1, u_2\}$  is a detour basis, weak edge detour basis and an edge detour basis for  $G$  so that  $dn(G) = dn_w(G) = dn_1(G) = 2$ . The graphs  $G$  given in Figure 2 are edge detour graphs.

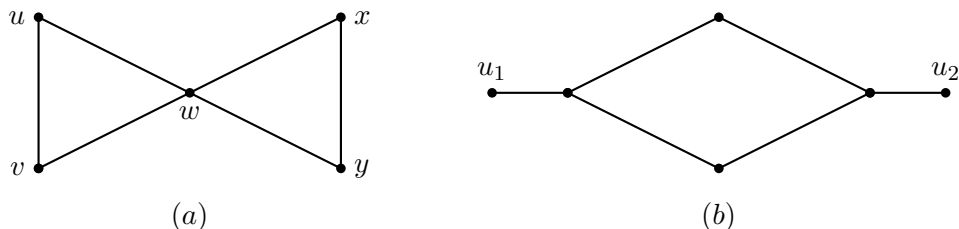


FIGURE 2.  $G$

For the graph  $G$  given in Figure 1, the set  $S = \{v_1, v_2\}$  is a detour basis and also a weak edge detour basis. However, it does not contain an edge detour set and so the graph  $G$  in Figure 1 is not an edge detour graph.

The following theorems are used in the sequel.

**Theorem 1** ([11]). *For any edge detour graph  $G$  of order  $p \geq 2$ ,  $2 \leq dn_1(G) \leq p$ .*

**Theorem 2** ([11]). *If  $G$  is an edge detour graph of order  $p \geq 3$  such that  $\{u, v\}$  is an edge detour basis of  $G$ , then  $u$  and  $v$  are not adjacent.*

A vertex of degree 1 is an *end-vertex* of  $G$ .

**Theorem 3** ([11]). *Every end-vertex of an edge detour graph  $G$  belongs to every edge detour set of  $G$ . Also if the set  $S$  of all end-vertices of  $G$  is an edge detour set, then  $S$  is the unique edge detour basis for  $G$ .*

**Theorem 4** ([11]). *If  $T$  is a tree with  $k \geq 2$  end-vertices, then  $dn_1(T) = k$ .*

**Theorem 5** ([11]). *Let  $G$  be an edge detour graph with cut-vertices and  $S$  an edge detour set of  $G$ . Then for any cut-vertex  $v$  of  $G$ , every component of  $G - v$  contains an element of  $S$ .*

Throughout this paper  $G$  denotes a connected graph with at least two vertices.

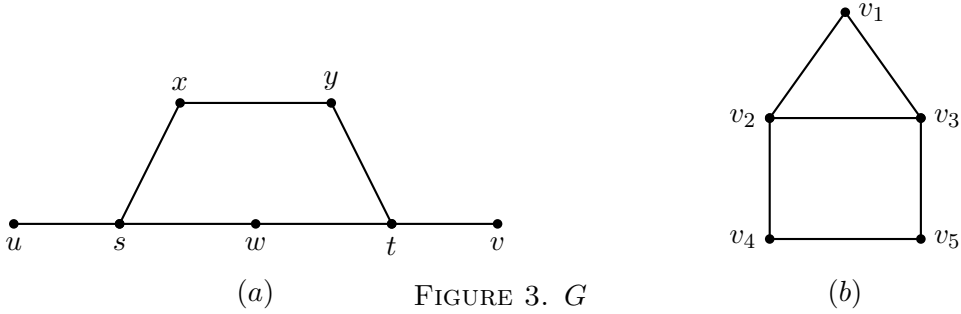
## 2. FORCING EDGE DETOUR SETS IN A GRAPH

Forcing geodetic number and forcing dimension of a graph was introduced and studied in [5, 6].

In this section we introduce the forcing edge detour number of an edge detour graph and determine the same for some standard classes of edge detour graphs.

**Definition 1.** Let  $G$  be an edge detour graph and  $S$  an edge detour basis of  $G$ . A subset  $T \subseteq S$  is called a forcing subset for  $S$  if  $S$  is the unique edge detour basis containing  $T$ . A forcing subset for  $S$  of minimum cardinality is a minimum forcing subset of  $S$ . The forcing edge detour number of  $S$ , denoted by  $fdn_1(S)$ , is the cardinality of a minimum forcing subset for  $S$ . The forcing edge detour number of  $G$ , denoted by  $fdn_1(G)$ , is  $fdn_1(G) = \min \{fdn_1(S)\}$ , where the minimum is taken over all edge detour bases  $S$  in  $G$ .

**Example 1.** For the graph  $G$  given in Figure 2(a),  $S = \{u, v, x, y\}$  is the unique edge detour basis of  $G$  so that  $fdn_1(G) = 0$ . For the graph  $G$  given in Figure 3 (a),  $S_1 = \{u, v, x\}$ ,  $S_2 = \{u, v, y\}$  and  $S_3 = \{u, v, w\}$  are the only edge detour bases of  $G$  and so  $fdn_1(G) = 1$ . For the graph  $G$  given in Figure 3 (b),  $S_1 = \{v_1, v_2, v_4\}$ ,  $S_2 = \{v_1, v_2, v_5\}$ ,  $S_3 = \{v_1, v_3, v_4\}$ ,  $S_4 = \{v_1, v_3, v_5\}$  and  $S_5 = \{v_1, v_4, v_5\}$  are the only five edge detour bases of  $G$  so that  $fdn_1(G) = 2$ .



The next three theorems are easy consequences of the respective definitions.

**Theorem 6.** For every edge detour graph  $G$ ,  $0 \leq fdn_1(G) \leq dn_1(G)$ .

**Remark 1.** The bounds in Theorem 6 are sharp. For the graph  $G$  given in Figure 2(a),  $fdn_1(G) = 0$ . For the complete graph  $K_4$ ,  $fdn_1(K_4) = dn_1(K_4) = 3$ . Also, all the inequalities in Theorem 6 can be strict. For the graph  $G$  given in Figure 3(a),  $fdn_1(G) = 1$  and  $dn_1(G) = 3$  so that  $0 < fdn_1(G) < dn_1(G)$ .

**Theorem 7.** Let  $G$  be an edge detour graph. Then

- i)  $fdn_1(G) = 0$  if and only if  $G$  has a unique edge detour basis,
- ii)  $fdn_1(G) = 1$  if and only if  $G$  has at least two edge detour bases, one of which is a unique edge detour basis containing one of its elements, and

iii)  $fdn_1(G) = dn_1(G)$  if and only if no edge detour basis of  $G$  is the unique edge detour basis containing any of its proper subsets.

**Theorem 8.** *Let  $G$  be an edge detour graph and  $W$  the set of all edge detour vertices of  $G$ . Then  $fdn_1(G) \leq dn_1(G) - |W|$ .*

**Remark 2.** *The bound in Theorem 8 is sharp. For the graph  $G$  given in Figure 3(a),  $dn_1(G) = 3$ ,  $|W| = 2$  and  $fdn_1(G) = 1$  as in Example 1. Also, the inequality in Theorem 8 can be strict. For the cycle  $C_4$ ,  $dn_1(C_4) = 2$ ,  $|W| = 0$  and  $fdn_1(C_4) = 1$ . Thus  $fdn_1(G) < dn_1(G) - |W|$ .*

The following theorems give the forcing edge detour numbers of certain classes of graphs.

**Theorem 9.** *Let  $G$  be a complete graph  $K_p$  or an odd cycle  $C_p$  of order  $p \geq 3$ . Then a set  $S \subseteq V$  is an edge detour basis of  $G$  if and only if  $S$  consists of any three vertices of  $G$ .*

*Proof.* Let  $G = K_p$  ( $p \geq 3$ ). If  $\{u, v\}$  is any set of two vertices of  $G$ , then all the edges of  $G$  other than  $uv$  lie on a  $u-v$  detour. Hence it follows that no two element subset of  $V$  is an edge detour set of  $G$ . Let  $S = \{u, v, w\}$  be any set of three vertices of  $G$ . Then every edge  $e$  of  $G$  lies on a detour joining a pair of vertices of  $S$ . Hence it follows that  $S$  is an edge detour basis of  $G$ . Now, assume that  $S$  is an edge detour basis of  $G$ . By Theorem 2,  $|S| \geq 3$ . It follows from the first part of the proof that  $|S| = 3$  and  $S$  consists of any three vertices of  $G$ .

Let  $G$  be an odd cycle  $C_p$  ( $p \geq 3$ ). If  $\{u, v\}$  is any set of two vertices of  $G$ , then no edge on the  $u-v$  geodesic lies on the  $u-v$  detour in  $G$  and so no two element subset of  $V$  is an edge detour set of  $G$ . Let  $S = \{u, v, w\}$  be any set of three vertices of  $G$ . Then every edge in  $G$  lies on any one of the  $u-v$ ,  $v-w$  or  $u-w$  detours so that  $S$  is an edge detour basis of  $G$ . Now, assume that  $S$  is an edge detour basis of  $G$ . Since any set of two vertices of  $G$  is not an edge detour set of  $G$ , it follows just as above that  $|S| = 3$  and  $S$  consists of any three vertices of  $G$ .  $\square$

If  $u$  and  $v$  are two vertices in a graph  $G$ , then the *distance*  $d(u, v)$  is the length of a shortest  $u-v$  path in  $G$ . A  $u-v$  path of length  $d(u, v)$  is a  $u-v$  *geodesic*. The *diameter*  $d(G)$  of a connected graph  $G$  is the length of any longest geodesic. Two vertices  $u, v$  in  $G$  are *antipodal* if  $d(u, v) = d(G)$ .

**Theorem 10.** *Let  $G$  be an even cycle. Then a set  $S \subseteq V$  is an edge detour basis of  $G$  if and only if  $S$  consists of two antipodal vertices of  $G$ .*

*Proof.* It is clear that every set  $S$  of two antipodal vertices of  $G$  is an edge detour set and so  $S$  is an edge detour basis of  $G$ . On the other hand, if  $u$  and  $v$  are not antipodal vertices, then the edges of  $u-v$  geodesic do not lie

on the  $u$ - $v$  detour in  $G$  and so the set  $\{u, v\}$  is not an edge detour set of  $G$ . Therefore, every edge detour basis of  $G$  must consist of two antipodal vertices of  $G$ .  $\square$

**Corollary 11.** *For a complete graph  $K_p$  ( $p \geq 3$ ),*

- i)  $dn_1(K_3) = 3$  and  $fdn_1(K_3) = 0$ .
- ii)  $dn_1(K_p) = fdn_1(K_p) = 3$  for  $p \geq 4$ .

*Proof.* i) By Theorem 9,  $dn_1(K_3) = 3$ . Let  $v_1, v_2, v_3$  be the vertices of  $K_3$ . Then, by Theorem 9, it is clear that the set  $\{v_1, v_2, v_3\}$  is the unique edge detour basis of  $K_3$  and hence by Theorem 7(i),  $fdn_1(K_3) = 0$ .

ii) Let  $p \geq 4$ . By Theorem 9,  $dn_1(K_p) = 3$ . Since  $p \geq 4$ , it follows from Theorem 9 that no subset of  $V$  of cardinality at most 2 is a forcing subset for any edge detour basis of  $K_p$ . Therefore, by Theorem 7(iii),  $fdn_1(K_p) = dn_1(K_p) = 3$ .  $\square$

**Corollary 12.** *If  $G$  is the cycle  $C_p$  ( $p \geq 4$ ), then*

- i)  $dn_1(G) = 2$  and  $fdn_1(G) = 1$  for  $p$  even.
- ii)  $dn_1(G) = fdn_1(G) = 3$  for  $p$  odd.

*Proof.* i) Let  $p$  be even. It follows from Theorem 10 that  $dn_1(G) = 2$  and that each vertex in  $G$  belongs to exactly only one edge detour basis of  $G$ . Hence every set consisting of a single vertex of  $G$  is a forcing subset for an edge detour basis of  $G$  so that  $fdn_1(G) = 1$ .

ii) Let  $p$  be odd. Then the result follows from Theorems 9 and 7(iii) and the proof is similar to that of Corollary 11(ii).  $\square$

A set  $S$  of vertices in a graph is *independent* if no two vertices in  $S$  are adjacent.

**Theorem 13.** *Let  $G$  be a complete bipartite graph  $K_{m,n}$  ( $2 \leq m \leq n$ ). Then a set  $S \subseteq V$  is an edge detour basis of  $G$  if and only if  $S$  consists of any two independent vertices of  $G$ .*

*Proof.* Let  $X$  and  $Y$  be the partite sets of  $G$  with  $|X| = m$  and  $|Y| = n$ . Let  $S = \{u, v\}$  be any set of two independent vertices in  $G$ . Then it is clear that  $D(u, v) = 2m - 2$  or  $D(u, v) = 2m$  according to whether  $u, v \in X$  or  $u, v \in Y$ . First assume that  $u, v \in X$ . Let  $xy$  be an edge such that  $x \in X$  and  $y \in Y$ . If  $x \neq u$ , then the edge  $xy$  lies on the  $u$ - $v$  detour  $P : u, y, x, \dots, v$  of length  $2m - 2$ . If  $x = u$ , then the edge  $xy$  lies on the  $u$ - $v$  detour  $P : u = x, y, \dots, v$  of length  $2m - 2$ . Hence  $S$  is an edge detour set of  $G$ . The case when  $u, v \in Y$  is similar.

Now, assume that  $S$  is an edge detour basis of  $G$ . It follows from the first part of the proof that  $|S| = 2$ . Let  $S = \{u, v\}$ . Then, by Theorem 2,  $u$  and  $v$  are not adjacent. Thus  $S$  consists of two independent vertices of  $G$ .  $\square$

**Corollary 14.** For a complete bipartite graph  $G = K_{m,n}$  ( $2 \leq m \leq n$ ),

i)  $dn_1(G) = 2$  and  $fdn_1(G) = 1$  for  $m = 2$  and  $n \geq 2$ .

ii)  $dn_1(G) = fdn_1(G) = 2$  for  $m, n \geq 3$ .

*Proof.* Let  $X = \{u_1, u_2, \dots, u_m\}$  and  $Y = \{v_1, v_2, \dots, v_n\}$  be the bipartite sets of  $K_{m,n}$  ( $2 \leq m \leq n$ ).

i) If  $m = 2$  and  $n = 2$ , then  $G = C_4$  and the result follows from Corollary 12(i). Let  $m = 2$  and  $n \geq 3$ . By Theorem 13,  $dn_1(G) = 2$ . By Theorem 13, the set  $X = \{u_1, u_2\}$  is the unique edge detour basis of  $G$  such that  $\{u_1\}$  is a forcing subset for  $X$ . Since  $n \geq 3$ , there is more than one edge detour basis and it follows that  $fdn_1(G) = 1$ .

ii) By Theorem 13,  $dn_1(G) = 2$ . Also it follows from Theorem 13 that each vertex belongs to more than one edge detour basis of  $G$  and so  $fdn_1(G) > 1$ . Since  $dn_1(G) = 2$ , it follows that  $fdn_1(G) = 2$ .  $\square$

**Theorem 15.** If  $T$  is a tree with  $k \geq 2$  end-vertices, then  $dn_1(G) = k$  and  $fdn_1(G) = 0$ .

*Proof.* By Theorem 4,  $dn_1(G) = k$ . Since the set of all end-vertices of a tree is the unique edge detour basis, the result follows from Theorem 7(i).  $\square$

In view of Theorem 6, we have the following realization result.

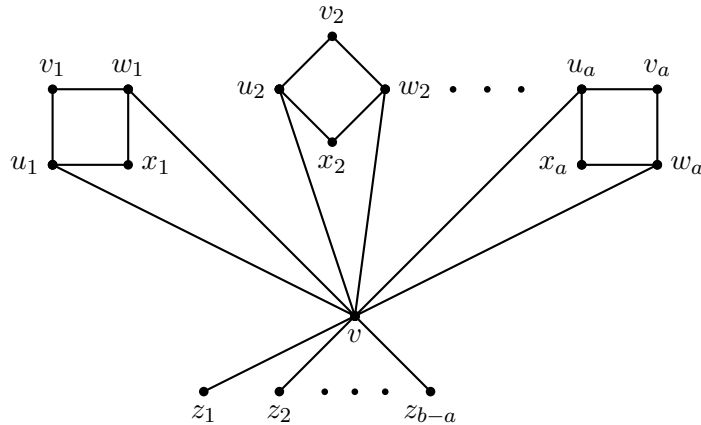
**Theorem 16.** For each pair  $a, b$  of integers with  $0 \leq a \leq b$  and  $b \geq 2$ , there is an edge detour graph  $G$  with  $fdn_1(G) = a$  and  $dn_1(G) = b$ .

*Proof. Case 1:*  $a = 0$ . For each  $b \geq 2$ , let  $G$  be a tree with  $b$  end-vertices. Then  $fdn_1(G) = 0$  and  $dn_1(G) = b$  by Theorem 15.

**Case 2:**  $a \geq 1$ . For each  $i$  ( $1 \leq i \leq a$ ), let  $F_i : u_i, v_i, w_i, x_i, u_i$  be the cycle of order 4 and let  $H = K_{1,b-a}$  be the star at  $v$  whose set of end-vertices is  $\{z_1, z_2, \dots, z_{b-a}\}$ . Let  $G$  be the graph obtained by joining the central vertex  $v$  of  $H$  to the vertices  $u_i$  and  $w_i$  of each  $F_i$  ( $1 \leq i \leq a$ ). Clearly the graph  $G$  is connected and is shown in Figure 4.

Let  $W = \{z_1, z_2, \dots, z_{b-a}\}$  be the set of all  $(b-a)$  end-vertices of  $G$ . First, we show that  $dn_1(G) = b$ . By Theorems 3 and 5, every edge detour basis contains  $W$  and at least one vertex from each  $F_i$  ( $1 \leq i \leq a$ ). Thus  $dn_1(G) \geq (b-a) + a = b$ . On the other hand, since the set  $S_1 = W \cup \{v_1, v_2, \dots, v_a\}$  is an edge detour set of  $G$ , it follows that  $dn_1(G) \leq |S_1| = b$ . Therefore  $G$  is an edge detour graph and  $dn_1(G) = b$ .

Next we show that  $fdn_1(G) = a$ . It is clear that  $W$  is the set of all edge detour vertices of  $G$ . Hence it follows from Theorem 8 that  $fdn_1(G) \leq dn_1(G) - |W| = b - (b-a) = a$ . Now, since  $dn_1(G) = b$ , it is easily seen that a set  $S$  is an edge detour basis of  $G$  if and only if  $S$  is of the form  $S = W \cup \{y_1, y_2, \dots, y_a\}$ , where  $y_i \in \{v_i, x_i\} \subseteq V(F_i)$  ( $1 \leq i \leq a$ ). Let  $T$  be a subset of  $S$  with  $|T| < a$ . Then there is a vertex  $y_j$  ( $1 \leq j \leq a$ ) such that

FIGURE 4.  $G$ 

$y_j \notin T$ . Let  $s_j \in \{v_j, x_j\} \subseteq V(F_j)$  distinct from  $y_j$ . Then  $S' = (S - \{y_j\}) \cup \{s_j\}$  is an edge detour basis that contains  $T$ . Thus  $S$  is not the unique edge detour basis containing  $T$ . Thus  $fdn_1(S) \geq a$ . Since this is true for all edge detour bases of  $G$ , it follows that  $fdn_1(G) \geq a$  and so  $fdn_1(G) = a$ .  $\square$

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